

CLIP OS: Building a defense-in-depth OS around Linux kernel security improvements

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About the ANSSI

- Agence nationale de la sécurité des systèmes d'information
- French authority in the area of cyberdefence, network and information security
- We are **not** an intelligence agency

Overview

CLIP OS ?

- Linux distribution developed by the ANSSI
- Initially only available internally
- ▶ Now open source, mostly under the LGPL v2.1+
- Code and issue tracker hosted on GitHub:
 - Version 4: available as reference and for upstream patch contribution¹
 - Version 5: currently developed version, alpha status²

¹https://github.com/CLIPOS-Archive ²https://github.com/CLIPOS



- Hardened Linux kernel and userspace
- Confined services
- "Unprivileged" admin, audit and update roles:

 the root account is not usable
- Automatic updates using A/B partition model (similar to Android 7+)

Multilevel security OS

Provide two isolated user environments: low and high

- Interactions follow the Bell-LaPadula model:
 - Write up: upload documents from *low* to *high*
 - Read down: high has read only access to untrusted USB devices
 - Trusted write down: encrypt documents from high to write them in low
- Level high can only access network through a VPN
- Per level user device assignment

Multilevel from the end user point of view



Admin panel: devices assignment per level

×	0	Socie: Attribution des périphériques	Bureau
0	CD-ROM	Périphériques d'entrée/sortie	
×	Dossier personnel	Attribution de la carte son (immédiat) :	AL
B to	Support USB	Attribution de l'imprimante USB (prochain branchement)	:
		Attribution du scanner (prochain branchement) :	
20	Corbeille	Attribution de la webcam (prochain branchement) :	
26		Uuitter	X (1) + 20:26

Differences with Qubes OS

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Technical point of view

- Hypervisor vs. supervisor isolation
- Limited access right, even for the administrator

Architecture





Gentoo Hardened

- Hardened toolchain
- Flexible patching

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- Kernel self-protection (e.g., memory protection, CFI)
- Multiple userspace hardening features (e.g., chroot, TPE)

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CLIP LSM

- Complement the Linux permission model
- Leverage Linux-VServer and grsecurity/PaX

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The O_MAYEXEC flag

Enforce and extend $W \oplus X$ from mount points to scripts (via interpreters)





Partitioning

Hardened containers

- Leverage Linux-VServer admin and watch (audit) concepts
- New capability bounding sets: for root and per container
- Hardened chroot

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Container content and interaction

- Tailored filesystem layouts per service
- Container management with vsctl and clip-libvserver (self-jailing)

Veriexec and permissions (CLIP-LSM)

Goal

- Split Linux capabilities (e.g., Fuse, unshare)
- Add new permissions (e.g., network, XFRM)
- Can be tied to an XID
- Does not use xattr (thus independent from the filesystem)

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Configuration example: /etc/verictl.d/chromium

/usr/.../chrome-sandbox 1002 e
SETUID|SETGID|SYS_CHROOT SETUID|SETGID|SYS_CHROOT cUP sha256 45bcbd1...

Veriexec example





General Linux kernel hardening

Strict whitelist of kernel options, but easily composable sets

- Paranoid command line
 - iommu=force, pti=on, spectre_v2=on, etc.
- Strict sysctl defaults
 - kernel.kptr_restrict, kernel.yama.ptrace_scope, etc.

Enabling Linux kernel hardening

Goals

- Protecting the kernel from itself and from userspace
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Interaction with upstream & KSPP

- Include in-progress or ready-for-upstream patches
- Integrate and validate patches in a single tree
- Maintain hardening patches for latest stable kernel

Patch series: linux-hardened

- Memory hardening improvements, including:
 - better userspace ASLR
 - slab allocators hardening (mostly SLUB)
 - simpler page sanitizing
- Various restrictions: TIOCSTI ioctl, perf subsystem, device timing side channels, etc.
- Miscellaneous additions: more BUG_ONs, more ___ro_after_init, etc.

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- Development status: In progress
- CLIP OS status: Merged
- Upstream status: Most changes unlikely to be merged upstream

Upstream contribution integration: Lockdown

Features

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- Reduce information leaks and block attacks using uninitialized kernel stack variables:
 - Erase the stack before returning from system calls
- Improve runtime detection of kernel stack overflows (e.g. Stack Clash):
 - Instrument calls to alloca()

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- Stackable LSM
- Powered by eBPF
- Dynamic filesystem access control using whitelists & blacklists
- See landlock.io

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- Stackable LSM
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- Dynamic filesystem access control using whitelists & blacklists
- See landlock.io
- Development status: Initial feature set ready
- CLIP OS status: Planned
- Upstream status: Work in progress

Upstream contribution: VServer-like LSM

- Adds a single kernel enforced indentifier for confined environments
- Similar in principle to VServer XID or to "Container IDs"
- Inspired by the VServer patch
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- Development status: Early development stage
- CLIP OS status: Planned

Conclusion

Take away

- Hardened Linux distro and kernel
- Coordinated userspace and kernelspace
- Support multilevel security

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Ongoing project

- Contributions welcome
- Browse the doc and the sources to find more interesting features: docs.clip-os.org

Thanks!

Sclip-os.org

⊠ clipos@ssi.gouv.fr

• v4: github.com/CLIPOS-Archive

• v5: github.com/CLIPOS

We're hiring! (but not directly for CLIP OS) Linux system security expert https://www.ssi.gouv.fr/emploi/expert-en-securite-des-systemes-linux/

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 - Linux kernel
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 - kernel command line

	UEFI Firmware
4	Bootloader
$\left(\right)$	EEL Binary
	(Linux kernel + initramfs + kernel command line)

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- Requires enrollment in hardware
- 2 Minimal bootloader (gummiboot/systemd-boot)
- 3 EFI bundle:
 - Linux kernel
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- 4 DM-Verity partition:
 - DM-Verity root hash set in kernel command line
 - Forward error correction support (FEC)
 - Read only uncompressed SquashFS root filesystem

